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## Is a Quantum Logic a Logic?

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(1. V. 70)

In a recent study Jauch and Piron [2] have considered the possibility that a quantum proposition system is an infinite valued logic. They argue that if this is the case then for any two propositions  $p$  and  $q$  there must exist a conditional proposition  $p \rightarrow q$ . Following Lukasiewicz [3] the truth value  $[p \rightarrow q]$  of the conditional  $p \rightarrow q$  is defined as follows:  $[p \rightarrow q] = \min \{1, 1 - [p] + [q]\}$  where  $[p]$  and  $[q]$  are the truth values of  $p$  and  $q$  respectively. Here  $[p] = 1$  is interpreted as ' $p$  is true'. Note that  $[p] = 1$  and  $[p \rightarrow q] = 1$  implies  $[q] = 1$  so we have a law of deduction, which is a property that any reasonable logic should possess. Notice further that if  $[p \rightarrow q] = 1$  and  $[q \rightarrow r] = 1$  then  $[p \rightarrow r] = 1$  so that implication is transitive as it should be.

Let  $\mathcal{L}$  be an orthomodular poset (representing some quantum proposition system) and let  $\mathcal{S}$  be an order determining (full in [1]) set of states on  $\mathcal{L}$ . We further assume that if  $m_1, m_2 \in \mathcal{S}$ , then  $1/2 m_1 + 1/2 m_2 \in \mathcal{S}$ , that is,  $\mathcal{S}$  is closed under the formation of mid-points. We say that  $a, b \in \mathcal{L}$  are *conditional* if there exists  $c \in \mathcal{L}$  such that for all  $m \in \mathcal{S}$   $m(c) = \min \{1, m(a) + m(b)\}$ . If  $c$  exists it is unique. We call  $c$  the *conditional* of  $a$  and  $b$  and write  $c = a \rightarrow b$ . We say that  $\mathcal{L}$  (or, more correctly, the pair  $(\mathcal{L}, \mathcal{S})$ ) is *conditional* if every pair  $a, b \in \mathcal{L}$  are conditional. Now if  $\mathcal{L}$  is to be a logic with a law of deduction then  $\mathcal{L}$  must be conditional. Jauch and Piron [2] have shown that standard proposition systems (that is, ones that are isomorphic to the lattice of all closed subspaces of a Hilbert space) are not conditional and thus cannot be logics in the usual sense. We generalize their results to the orthomodular posets  $\mathcal{L}$  considered above. In fact we obtain the strong result that  $\mathcal{L}$  is conditional if and only if  $\mathcal{L} = \{0, 1\}$ . We then characterize the pairs  $a, b \in \mathcal{L}$  which are conditional.

Undefined terms appear in [1]. If  $a \leq b'$  we write  $a + b$  for  $a \vee b$ . If  $a \leq b$  we write  $b - a$  for  $b \wedge a'$ . We first state a useful lemma whose simple proof is left to the reader.

*Lemma 1.* (i)  $m(a \rightarrow b) = 1$  if and only if  $m(a) \leq m(b)$ ;  $m(a \rightarrow b) = m(a') + m(b)$  if and only if  $m(b) \leq m(a) = 1$ .

(ii)  $m(a \rightarrow b) = m(b)$  if and only if  $m(b) = 1$  or  $m(a) = 1$ .

This lemma will be frequently used without further comment.

*Theorem 2.*  $\mathcal{L}$  is conditional if and only if  $\mathcal{L} = \{0, 1\}$ .

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*Proof:* Clearly  $\{0, 1\}$  is conditional; in fact  $1 = 0 \rightarrow 1$  and  $0 = 1 \rightarrow 0$ . Now let  $\mathcal{L}$  be conditional and suppose there exists  $a \in \mathcal{L} - \{0, 1\}$ . Then  $c = a \rightarrow a'$  exists and  $m(c) = \min\{1, 2m(a')\}$ . Since  $\mathcal{S}$  is order determining  $a' \leq c$ . Hence there exists  $b \in \mathcal{L}$  such that  $a' + b = c$ . Now  $m(b) = m(c) - m(a') = \min\{m(a), m(a')\}$ . Thus  $m(b) \leq 1/2$  for all  $m \in \mathcal{S}$ . It follows that  $b \leq b'$  since  $\mathcal{S}$  is order determining. Hence  $b = 0$  and  $c = a'$ . Thus  $m(c) = \min\{1, 2m(c)\}$  and hence  $m(c) = 0$  or  $1$  for all  $m \in \mathcal{S}$ . Moreover since  $0 < c < 1$  there exist  $m_1, m_2 \in \mathcal{S}$  with  $m_1(c) = 0$  and  $m_2(c) = 1$ . Letting  $m = 1/2 m_1 + 1/2 m_2$  we have  $m(c) = 1/2$ , a contradiction. Hence  $\mathcal{L} = \{0, 1\}$ .

We have seen that, for non-trivial posets  $\mathcal{L}$ , not every pair of elements is conditional. We now study the properties of pairs of elements that are conditional.

*Lemma 3.* If  $a \rightarrow b$  and  $a' \vee b$  exist and are equal then  $a C b$ .

*Proof:* There exists  $d \in \mathcal{L}$  such that  $b + d = a' \vee b$ . We show  $d \leq a'$ . Otherwise there exists  $m \in \mathcal{S}$  such that  $m(d) > m(a')$ . Then  $m(a' \vee b) = m(b) + m(d) > 1 - m(a) + m(b)$  so  $m(a) > m(b)$ . Hence  $m(a' \vee b) = m(a \rightarrow b) = 1 - m(a) + m(b)$ , a contradiction. Now there exists  $e \in \mathcal{L}$  with  $d + e = a'$ . We show  $e \leq b$ . Otherwise there exists  $m \in \mathcal{S}$  with  $m(e) > m(b)$ . Then  $m(a') = m(d) + m(e) > m(d) + m(b) = m(a' \vee b) \geq m(a')$ , a contradiction. Hence there exists  $f \in \mathcal{L}$  with  $b = f + e$ ,  $a' = d + e$  and  $f \leq b \leq d'$  so that  $a' C b$ . Thus  $a C b$ .

*Lemma 4.* If  $c = a \rightarrow b$  exists then  $a' \leq c$  and  $b \leq c$ .

*Proof:* If  $a' \leq c$  then there exists  $m \in \mathcal{S}$  such that  $m(c) < m(a')$ . Hence  $m(c) < 1$  and  $1 - m(a) + m(b) = m(c) < 1 - m(a)$ . Thus  $m(b) < 0$ , a contradiction. That  $b \leq c$  is immediate.

We say that  $\mathcal{S}$  is *sufficient* if  $0 \neq a \in \mathcal{L}$  implies there exists  $m \in \mathcal{S}$  with  $m(a) = 1$ .

*Theorem 5.* Let  $\mathcal{S}$  be sufficient and assume that  $a' \vee b$  exists. Then  $a \rightarrow b$  exists if and only if  $a \leq b$  or  $b \leq a$ .

*Proof:* Clearly, if  $a \leq b$  then  $a \rightarrow b = 1$  and if  $b \leq a$ , then  $a \rightarrow b = a' + b$ . Conversely, assume  $c = a \rightarrow b$  exists. By Lemma 4  $c \geq a' \vee b$ . Hence there exists  $d \in \mathcal{L}$  such that  $(a' \vee b) + d = c$ . Suppose  $d \neq 0$ . Then there exists  $m \in \mathcal{S}$  such that  $m(d) = 1$ . Hence  $m(a') = m(b) = 0$  and  $m(c) = 1 - m(a) + m(b) = 0$ , a contradiction. Therefore  $d = 0$  and  $c = a' \vee b$ . It now follows from Lemma 3 that  $a C b$ . Suppose  $a$  and  $b$  are not comparable. Then  $a \wedge b < a$  and  $a \wedge b < b$ . Hence there exists  $m_1, m_2 \in \mathcal{S}$  such that  $m_1(a - (a \wedge b)) = 1$  and  $m_2(b - (a \wedge b)) = 1$ . It follows that  $m_1(a) = m_2(b) = 1$  and  $m_1(b) = m_1(a \wedge b) = m_2(a) = m_2(a \wedge b) = 0$ . Let  $m = 1/2 (1/2 m_1 + 1/2 m_2) + 1/2 m_1 = 3/4 m_1 + 1/4 m_2$ . Then  $m(a \wedge b) = 0$  and  $m(b) = 1/4 < 3/4 = m(a)$ . Hence  $m(a') + m(b) = m(c) = m(a' \vee b) = m(a' + (a \wedge b)) = m(a') + m(a \wedge b)$ . Thus  $m(b) = m(a \wedge b)$ , a contradiction.

*Corollary 6.* Let  $\mathcal{S}$  be sufficient and  $a' \vee b$  exist. If  $a \rightarrow b$  exists, then  $a \rightarrow b = a' \vee b$ ,  $b \rightarrow a$  exists,  $b' \vee a$  exists, and  $b \rightarrow a = b' \vee a$ .

The proofs of the previous theorems depend heavily on the fact that  $\mathcal{S}$  is order determining, sufficient or both. If we strengthen  $\mathcal{S}$  still further we obtain a stronger result. We say that  $\mathcal{S}$  is *strongly order determining* if  $\{m \in \mathcal{S} : m(a) = 1\} \subset \{m \in \mathcal{S} : m(b) = 1\}$

implies that  $a \leq b$ . It can be shown that strongly order determining implies both order determining and sufficiency. (The converse fails; see [1].) Notice that the set of states on the lattice of all closed subspaces of a Hilbert space is strongly order determining.

*Theorem 7.* If  $\mathcal{S}$  is strongly order determining, then  $a \rightarrow b$  exists if and only if  $a \leq b$  or  $b \leq a$ .

*Proof:* As in Theorem 5, if  $a$  and  $b$  are comparable, then  $a \rightarrow b$  exists. Now assume  $c = a \rightarrow b$  exists. Suppose  $a \not\leq b$  and  $b \not\leq a$ . Then there exists  $m_0, m_1 \in \mathcal{S}$  such that  $m_0(a) = 1$ ,  $m_0(b) < 1$ ,  $m_1(a) < 1$  and  $m_1(b) = 1$ . Note that  $m_0(c) = m_0(b)$  and  $m_1(c) = 1$ . Let  $m = 1/2 m_0 + 1/2 m_1$ . Then  $m(a) = 1/2 + 1/2 m_1(a) < 1$ ,  $m(b) = 1/2 m_0(b) + 1/2 < 1$  and  $m(c) = m(b)$ . This last sentence contradicts Lemma 1 (ii). Hence  $a$  and  $b$  are comparable.

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