

# Pull-backs and product tests

Autor(en): **Wilce, Alexander**

Objektyp: **Article**

Zeitschrift: **Helvetica Physica Acta**

Band (Jahr): **70 (1997)**

Heft 6

PDF erstellt am: **11.09.2024**

Persistenter Link: <https://doi.org/10.5169/seals-117053>

## **Nutzungsbedingungen**

Die ETH-Bibliothek ist Anbieterin der digitalisierten Zeitschriften. Sie besitzt keine Urheberrechte an den Inhalten der Zeitschriften. Die Rechte liegen in der Regel bei den Herausgebern.

Die auf der Plattform e-periodica veröffentlichten Dokumente stehen für nicht-kommerzielle Zwecke in Lehre und Forschung sowie für die private Nutzung frei zur Verfügung. Einzelne Dateien oder Ausdrucke aus diesem Angebot können zusammen mit diesen Nutzungsbedingungen und den korrekten Herkunftsbezeichnungen weitergegeben werden.

Das Veröffentlichen von Bildern in Print- und Online-Publikationen ist nur mit vorheriger Genehmigung der Rechteinhaber erlaubt. Die systematische Speicherung von Teilen des elektronischen Angebots auf anderen Servern bedarf ebenfalls des schriftlichen Einverständnisses der Rechteinhaber.

## **Haftungsausschluss**

Alle Angaben erfolgen ohne Gewähr für Vollständigkeit oder Richtigkeit. Es wird keine Haftung übernommen für Schäden durch die Verwendung von Informationen aus diesem Online-Angebot oder durch das Fehlen von Informationen. Dies gilt auch für Inhalte Dritter, die über dieses Angebot zugänglich sind.

## Pull-backs and Product Tests

By Alexander Wilce

Department of Mathematics and Computer Science  
Juniata College  
Huntingdon, PA 16652

(28.X.1996)

*Abstract.* Let  $\mathcal{A}$  and  $\mathcal{B}$  be test spaces. We study the test space  $B(\mathcal{A}, \mathcal{B})$  consisting of graphs of bijections  $f : E \rightarrow F$  between tests  $E \in \mathcal{A}$  and  $F \in \mathcal{B}$ . Elements of  $B(\mathcal{A}, \mathcal{B})$  may be interpreted as *products*, in something like the sense of Piron, of tests in  $\mathcal{A}$  and  $\mathcal{B}$ .

### Introduction

In a long series of papers (cf [2], [3], [4] and references therein), D. J. Foulis and the late C.H. Randall developed a straightforward but versatile generalized probability theory based on what are now usually called *test spaces*. In brief: A test space  $\mathcal{A}$  is simply a non-empty collection of discrete sets  $E, F, \dots$ , each thought of as the outcome-set for some measurement or *test*. When  $\mathcal{A}$  contains only one test, one recovers (discrete) classical probability theory; when it consists of the set of maximal orthonormal bases of a Hilbert space, one recovers quantum probability theory.

This note concerns the following construction: If  $\mathcal{A}$  and  $\mathcal{B}$  are test spaces, let  $B(\mathcal{A}, \mathcal{B})$  denote the set of bijections  $f : E \rightarrow F$  between tests  $E \in \mathcal{A}$  and  $F \in \mathcal{B}$ . Identifying each such a bijection with its graph,  $B(\mathcal{A}, \mathcal{B})$  may be regarded as a test space in its own right.

We propose to interpret  $B(\mathcal{A}, \mathcal{B})$  as the test space consisting of *products*, in something close to the sense of Piron [8] and Aerts [1], of tests  $E \in \mathcal{A}$  and  $F \in \mathcal{B}$ . The construction is also of interest on purely mathematical grounds. On the one hand, it preserves various standard regularity conditions on  $\mathcal{A}$  and  $\mathcal{B}$ ; on the other hand, as soon as  $\mathcal{A}$  and  $\mathcal{B}$  contain tests with more than two outcomes, the structure of  $B(\mathcal{A}, \mathcal{B})$  becomes quite rich, *even if  $\mathcal{A}$  and  $\mathcal{B}$  are classical*. Moreover, for certain categories of “uniform” test spaces,  $B(\mathcal{A}, \mathcal{B})$  is effective as the direct product of  $\mathcal{A}$  and  $\mathcal{B}$ .

In section 1, we discuss our construction in general terms. In section 2, we discuss the stability of various regularity conditions on  $\mathcal{A}$  and  $\mathcal{B}$  under passage to  $B(\mathcal{A}, \mathcal{B})$ . In particular, we show that if  $\mathcal{A}$  and  $\mathcal{B}$  are algebraic, then  $B(\mathcal{A}, \mathcal{B})$  is algebraic as well. In section 3, we characterize the logic of  $B(\mathcal{A}, \mathcal{B})$  in the case that  $\mathcal{A}$  and  $\mathcal{B}$  are algebraic.

## 1 Questions, Products and Pull-Backs

As explained above, a **test space**<sup>1</sup> is a non-empty set  $\mathcal{A}$  of non-empty sets  $E, F, \dots$ . Elements of  $\mathcal{A}$  are called **tests** and elements of  $X := \bigcup \mathcal{A}$  are called **outcomes**. The intended interpretation is that each test  $E \in \mathcal{A}$  is an exhaustive set of mutually exclusive outcomes, as, for instance, the set of outcomes of some experiment. Borrowing terminology from classical probability theory, we refer to any subset of any test  $E \in \mathcal{A}$  as an **event** of  $\mathcal{A}$ . We write  $\mathcal{E}(\mathcal{A})$  for the set of all events of  $\mathcal{A}$ .

Test spaces provide the foundation for a very natural – and conceptually uncomplicated – generalization of elementary probability theory having both classical measure-theoretic and quantum-mechanical probability as special cases. It is worth a moment to give a sketch of this. One defines a **state** on a test space  $\mathcal{A}$  to be a map  $\omega : X \rightarrow [0, 1]$  such that  $\omega(x) \geq 0$  for each  $x \in X$  and  $\sum_{x \in E} \omega(x) = 1$  for each test  $E \in \mathcal{A}$ . In other words, a state is a real-valued function on the set of outcomes that restricts to a probability weight on each test.

Note that if  $\mathcal{A}$  consists of but a single test – i.e., if  $\mathcal{A} = \{E\}$  — then a state is simply a discrete probability distribution and we recover discrete classical probability theory. In this case, we call  $\mathcal{A}$  a *classical test space*. One can also consider the test space consisting of all countable partitions of a measurable space by measurable sets; this may be called a *Kolmogorov test space*. A *quantum test space* (or *frame manual*) is the set  $\mathcal{A}$  of all orthonormal bases of a Hilbert space  $\mathbf{H}$ . The outcomes of  $\mathcal{A}$  are the unit vectors of  $\mathbf{H}$ . Gleason's theorem [5] allows us to identify the states  $\omega$  on  $\mathcal{A}$  with density operators  $W$  on  $\mathbf{H}$  via the prescription  $\omega(x) = \langle Wx, x \rangle$  (where  $x$  is a unit vector of  $\mathbf{H}$ ).

We may wish to attach numerical or other labels to the outcomes of a test. This motivates the following terminology:

**1.1 Definition:** Given a set  $V$ , we define a  **$V$ -valued question** on a test space  $\mathcal{A}$  to be a bijection<sup>2</sup>  $\alpha : E \rightarrow V$ , where  $E$  is a test belonging to  $\mathcal{A}$ . The question is *posed* by executing the test  $E$ ; its *answer* is the value  $\alpha(x) \in V$  corresponding to the secured outcomes  $x \in E$ .

Note that if  $V = \{\text{yes}, \text{no}\}$ , this corresponds to the notion of a question as defined in the work of Piron [8].

If  $\alpha : E \rightarrow V$  and  $\beta : F \rightarrow V$  are two  $V$ -valued questions, it is very natural to form their

<sup>1</sup>called also a *manual* or *generalized sample space* in the older literature

<sup>2</sup>the condition that  $\alpha$  be bijective is benign: If not, replace  $V$  by the range of  $\alpha$  and  $E$ , by the partition  $\{\alpha^{-1}(x) \mid x \in V\}$ .

pull-back — that is, the canonical bijection  $\alpha \cdot \beta : E \times_V F \rightarrow V$  where

$$E \times_V F = \{ (x, y) \in E \times F \mid \alpha(x) = \beta(y) \}.$$

$$\begin{array}{ccc} E \times_V F & \longrightarrow & F \\ \downarrow & & \downarrow \beta \\ E & \xrightarrow{\alpha} & V \end{array}$$

More generally, given an arbitrary collection  $\{\alpha_i\}_{i \in I}$  of  $V$ -valued questions  $\alpha_i : E_i \rightarrow V$ , one can construct  $E = \{ x \in \prod_{i \in I} E_i \mid \alpha_i(x_i) = \alpha_j(x_j) \ \forall i, j \in I \}$  and set  $\prod_i \alpha_i(x) = \prod_j \alpha_j$  for any  $j \in I$ . (Indeed, by iterating this construction and taking a suitable direct limit, one can construct a test space that is in some sense closed under the formation of products of  $V$ -valued questions. We shall not pursue this here.)

We may interpret  $E \times_V F$  as a test, as follows: One of the tests  $E$  or  $F$  is selected. If the outcome of the selected test is, say,  $x \in E$ , then the outcome of  $E \times_V F$  is the unique pair  $(x, y) \in E \times_V F$  having  $x$  as its first component. Similarly, if the secured outcome is  $y \in F$ , the outcome of  $E \times_V F$  is the unique pair  $(x, y)$  with  $y$  as its second component. (Note that this in effect erases any record of which of the tests  $E$  and  $F$  was in fact selected.) To pose the question  $\alpha \cdot \beta$ , one executes  $E \times_V F$ . Upon securing, say,  $(x, y)$ , one records the value  $\alpha(x) = \beta(y)$  as answer.

As the reader familiar with [8] will have recognized, this construction is analogous to the notion of a product of yes-no questions as defined by Piron:

If  $\{\alpha_i\}$  is a family of questions, we denote by  $\prod_i \alpha_i$  the question defined in the following manner: One measures an arbitrary one of the  $\alpha_i$  and attributes to  $\prod_i \alpha_i$  the answer thus obtained. ([8], p. 20).

This notion makes equal sense for  $V$ -valued questions generally, and we believe our construction adequately captures it in a precise way.

The balance of this paper is devoted to a discussion of the test space consisting of tests  $E \times_V F$  arising from the formation of products of  $V$ -valued questions. This turns out to have a surprisingly rich structure. Before carrying on, it will be helpful to reformulate the definition of  $E \times_V F$  in a manner not depending explicitly upon the questions  $\alpha$  and  $\beta$ . To this end, notice that  $E \times_V F$  is simply the graph of the bijection  $\beta^{-1} \circ \alpha : E \rightarrow F$ . Conversely, given any pair of tests  $E, F \in \mathcal{A}$  and any bijection  $f : E \rightarrow F$ , we may understand  $f$  as a test corresponding to a product of  $V$ -valued questions defined on  $E$  and  $F$ , respectively. (To execute the test represented by  $f$ , one chooses  $E$  or  $F$ , executes it, and records the pair  $(x, f(x))$  or  $(f^{-1}(y), y)$  according as  $x \in E$  or  $y \in F$  is secured.)

**1.2 Definition:** For two sets  $E$  and  $F$ , we denote by  $B(E, F)$  the set of (graphs of) bijections  $f : E \rightarrow F$ , abbreviating  $B(E, E)$  to  $B(E)$ . For any two test-spaces  $\mathcal{A}, \mathcal{B}$ , we denote by  $B(\mathcal{A}, \mathcal{B})$  the collection of sets  $B(E, F)$  with  $E \in \mathcal{A}$  and  $F \in \mathcal{B}$ . We abbreviate  $B(\mathcal{A}, \mathcal{A})$  to  $B(\mathcal{A})$ .

Of course,  $B(\mathcal{A}, \mathcal{B})$  may be empty. On the other hand, as  $B(E, E) \neq \emptyset$ ,  $\mathcal{B}(\mathcal{A})$  is always rather large. Indeed, if  $\mathcal{A}$  is a totally finite test space having  $k$  operations each with  $n$  outcomes,  $B(\mathcal{A})$  has  $k^2n!$  operations (each with  $n$  outcomes). There is a natural embedding of  $\mathcal{A}$  in  $B(\mathcal{A})$ , namely, the diagonal map  $X \rightarrow X \times X$  given by  $x \mapsto (x, x)$ . This maps each test  $E \in \mathcal{A}$  to the corresponding identity function  $\text{Id}_E$ .

In general, the set of outcomes of  $B(\mathcal{A}, \mathcal{B})$  will be smaller than  $X \times Y$  (since, e.g., there may be outcomes in the former that belong only to tests with  $n$  outcomes, and outcomes of the latter belonging only to  $k$ -outcome tests with  $k \neq n$ ). In any case, if  $(x, y)$  and  $(u, v)$  belong to  $\cup B(\mathcal{A}, \mathcal{B})$ , we have  $(x, y) \perp (u, v) \Rightarrow x \perp u \ \& \ y \perp v$ .

We now consider some examples.

**1.3 Example:** Suppose  $\mathcal{A}$  is a collection of pair-wise disjoint two-element sets. Then

$$B(\mathcal{A}) = \{ \{(x, u), (y, v)\} \mid x \perp y, u \perp v \},$$

likewise a collection of pairwise-disjoint two-element sets. Notice that  $B(\mathcal{A})$  is naturally isomorphic to the set of pairs  $\{(\{x, u\}, \{y, v\}) \mid x \perp y, u \perp v\}$ , which is the model for the manual of product questions given by Foulis, Piron and Randall in [4].

Once we admit test spaces having operations with more than two outcomes, the structure of  $B(\mathcal{A})$  becomes quite involved. This is nicely illustrated even by the simplest example:

**1.4 Example:** Consider the hypergraph  $\mathcal{A} = \{E\}$  consisting of a single three-outcome experiment  $E = \{x, y, z\}$ . Then  $B(\mathcal{A}) = B(E)$  is isomorphic to the three-by-three “window” manual:

$$\begin{array}{ccccc} (x, x) & - - - & (y, z) & - - - & (z, y) \\ | & & | & & | \\ (y, y) & - - - & (z, x) & - - - & (x, z) \\ | & & | & & | \\ (z, z) & - - - & (x, y) & - - - & (y, x) \end{array}$$

As  $B(E)$  contains four-loops but no three-loops, its logic is an orthomodular poset, but not an orthomodular lattice ([7]). The state-space of  $B(E)$  is in effect the convex set of doubly stochastic  $3 \times 3$  matrices.

**1.5 Example:** Consider a Hilbert space  $\mathbf{H}$  (of any dimension, over any field) and let  $\mathcal{A}$  be the associated quantum test space, i.e., the set of all (un-ordered) orthonormal bases of  $\mathbf{H}$ . Every bijection  $f : E \rightarrow F$  between two bases  $E, F \in \mathcal{A}$  extends uniquely to a unitary operator on  $\mathbf{H}$ . If  $U$  is such an operator, its graph is a closed subspace of  $\mathbf{H} \times \mathbf{H}$ , and hence a Hilbert space in its own right. An orthonormal basis for  $U$  is simply the graph of  $U|_E$  for some  $E \in \mathcal{A}$ . Hence,  $B(\mathcal{A})$  is just the union over all unitaries  $U$ , of the frame manuals of the corresponding subspaces  $U \leq \mathbf{H} \times \mathbf{H}$ . (It is interesting to note that the set of graphs of unitaries on  $\mathbf{H}$  constitutes a *partial Hilbert space* in the sense of Gudder [6].)

We now consider a restricted class of test spaces for which the construction  $\mathcal{A}, \mathcal{B} \mapsto B(\mathcal{A}, \mathcal{B})$  behaves in a particularly satisfactory manner.

**1.6 Definition:** Let  $\kappa$  be any cardinal. A test space  $\mathcal{A}$  is  $\kappa$ -uniform iff every test  $E \in \mathcal{A}$  has cardinality  $\kappa$ .

If  $\mathcal{A}$  and  $\mathcal{B}$  are both  $\kappa$ -uniform, then  $Z = X \times Y$  and, in this case,  $(x, y) \perp (u, v)$  iff  $x \perp u$  and  $y \perp v$ . The class of  $\kappa$ -uniform test spaces is large enough to include both classical test spaces  $\mathcal{A} = \{E\}$  with  $\#(E) = \kappa$  and also the frame manual of any Hilbert space of dimension  $\kappa$ . Notice also that if  $\mathcal{A}$  and  $\mathcal{B}$  are both  $\kappa$ -uniform, then so also is  $B(\mathcal{A}, \mathcal{B})$ . In fact, as we shall now see,  $B(\mathcal{A}, \mathcal{B})$  serves as the direct product of uniform test spaces, provided we define our morphisms correctly.

**1.7 Definition:** By a **uniform map** between two test spaces  $\mathcal{A}$  and  $\mathcal{B}$  with outcome-sets  $X$  and  $Y$ , respectively, we mean a function  $\phi : X \rightarrow Y$  such that  $\phi(\mathcal{A}) \subseteq \mathcal{B}$  and  $x_1 \perp x_2 \Rightarrow \phi(x_1) \perp \phi(x_2)$  for all  $x_i \in X$ . (In the language of [4]: a uniform map is a *positive, outcome-preserving interpretation*.)

Note that if  $\phi$  is a uniform map, then  $\phi$  is locally bijective, in that for every  $E \in \mathcal{A}$ ,  $\phi|_E : E \rightarrow \phi(E) \in \mathcal{B}$  is a bijection.

**1.8 Theorem:** Let  $\mathcal{A}$  and  $\mathcal{B}$  be  $\kappa$ -uniform. Then  $B(\mathcal{A}, \mathcal{B})$  is the direct product of  $\mathcal{A}$  and  $\mathcal{B}$  in the category of uniform test spaces and uniform maps.

Proof: Let  $\mathcal{A}$  and  $\mathcal{B}$  be  $\kappa$ -uniform test spaces with  $\bigcup \mathcal{A} = X$  and  $\bigcup \mathcal{B} = Y$ . Note that  $B(\mathcal{A}, \mathcal{B})$  is again  $\kappa$ -uniform, and that  $\bigcup B(\mathcal{A}, \mathcal{B}) = X \times Y$ . Let  $\pi_1$  and  $\pi_2$  be the projections of  $X \times Y$  onto  $X$  and  $Y$ , respectively. If  $(x, y) \perp (u, v) \in \bigcup B(\mathcal{A}, \mathcal{B})$ , then  $x \perp u$  and  $y \perp v$ , so  $\pi_i(x, y) \perp \pi_i(u, v)$  for  $i = 1, 2$ . If  $f \in B(\mathcal{A}, \mathcal{B})$ , then  $\pi_1(f) = \text{dom}(f) \in \mathcal{A}$ ; similarly,  $\pi_2(f) = \text{ran}(f) \in \mathcal{B}$ . Thus, both projections are uniform maps. It now suffices to show that if  $\mathcal{C}$  is a  $\kappa$ -uniform test space with  $\bigcup \mathcal{C} = Z$ , and  $\phi : Z \rightarrow X$  and  $\psi : Z \rightarrow Y$  are uniform maps, then  $\phi \times \psi : Z \rightarrow (X \times Y)$  is an uniform map. If  $z \perp w$ , then  $\phi(z) \perp \phi(w)$  and  $\psi(z) \perp \psi(w)$ ; hence,  $(\phi \times \psi)(z) \perp (\phi \times \psi)(w)$ . Now suppose  $E \in \mathcal{C}$ . We must show that  $(\phi \times \psi)(E)$  belongs to  $B(\mathcal{A}, \mathcal{B})$ . Because  $\phi|_E$  is a bijection, we have

$$(\phi \times \psi)(E) = \{ (\phi(z), \psi(z)) \mid z \in E \} = \{ (x, \psi(\phi^{-1}(x))) \mid x \in \phi(E) \}.$$

That is,  $(\phi \times \psi)(E) = \psi \circ (\phi|_E)^{-1} : \phi(E) \rightarrow \psi(F)$ . Since  $\psi|_F$  is bijective, this last belongs to  $B(\mathcal{A}, \mathcal{B})$ .

## 2 The Structure of $B(\mathcal{A}, \mathcal{B})$

In this section, we establish (Theorems 2.2, 2.3 and 2.5) that passage from  $\mathcal{A}$  and  $\mathcal{B}$  to  $B(\mathcal{A}, \mathcal{B})$  preserves each of three standard conditions often imposed on test spaces: That of being algebraic, that of being coherent (though here we need an additional uniformity assumption), and that of being regular.

Throughout this section, let  $\mathcal{A}$  and  $\mathcal{B}$  be test spaces with outcome-sets  $X$  and  $Y$ , respectively. As noted above, the outcome-set of  $B(\mathcal{A}, \mathcal{B})$  is in general a *proper* (possibly

empty) subset of  $Z \subseteq X \times Y$ . An event for  $B(\mathcal{A}, \mathcal{B})$  is any subset of the graph of a bijection  $f : E \rightarrow F$  with  $E \in \mathcal{A}$  and  $F \in \mathcal{B}$ . Evidently, any such subset is the graph of a bijection between two events  $A \subseteq E$  and  $B \subseteq F$ . Thus,

$$\mathcal{E}(B(\mathcal{A}, \mathcal{B})) \subseteq B(\mathcal{E}(\mathcal{A}), \mathcal{E}(\mathcal{B})).$$

Again, the inclusion is generally proper – indeed, it is easy to see we have identity iff  $\mathcal{A}$  and  $\mathcal{B}$  are  $n$ -uniform for some finite  $n$ .

Events  $A$  and  $B$  of a test space  $\mathcal{A}$  are said to be *complementary* – the short-hand is  $A \subset B$  – iff  $A \cap B = \emptyset$  and  $A \cup B \in \mathcal{A}$ . If  $A$  and  $B$  are both complementary to a common third event, one says that  $A$  and  $B$  are *perspective*, writing  $A \sim B$ . A test space is a *algebraic* (in the older literature, a *manual*) iff, given any events  $A, B$  and  $C$ ,  $A \sim B$  and  $B \subset C$  imply  $A \subset C$ .

**2.1 Lemma:** Let  $f : A \rightarrow A'$  and  $g : B \rightarrow B'$  be bijections belonging to  $\mathcal{E}(B(\mathcal{A}, \mathcal{B}))$ . Then

- (1)  $f \subset g$  iff  $A \subset B$  and  $A' \subset B'$ .
- (2)  $f \sim g$  iff  $A \sim B$  and  $A' \sim B'$ .

Proof: Note that (2) is an immediate consequence of (1). To establish (1), suppose  $A \subset B$  and  $A' \subset B'$ . Then  $f \cap g = \emptyset$  and  $f \cup g \in B(A \cup B, A' \cup B') \subseteq B(\mathcal{A}, \mathcal{B})$ ; thus,  $f \subset g$ . Conversely, if  $f \subset g$ , then  $f \cap g = \emptyset$  and  $f \cup g \in B(E, F)$  for some  $E \in \mathcal{A}, F \in \mathcal{B}$ . But then  $A \cup B = E \in \mathcal{A}$  and, as  $f \cup g$  is again a bijection, we must have  $A \cap B = \emptyset$  – whence,  $A \subset B$ . Also,  $A' \cup B' = f(A) \cup g(B) = F \in \mathcal{B}$ , and, again because  $f \cup g$  is a bijection,  $A' \cap B' = \emptyset$ , so  $A' \subset B'$ .

**2.2 Theorem:** If  $\mathcal{A}$  and  $\mathcal{B}$  are algebraic, then  $B(\mathcal{A}, \mathcal{B})$  is likewise algebraic. If the test space  $B(\mathcal{A})$  is algebraic, then  $\mathcal{A}$  is algebraic.

Proof: Suppose that  $f : A \rightarrow A', g : B \rightarrow B'$  in  $\mathcal{E}(B(\mathcal{A}, \mathcal{B}))$  with  $f \sim g$  and  $g \subset h : C \rightarrow C'$ . By Lemma 1,  $A \sim B \subset C$  and  $A' \sim B' \subset C'$ . If  $\mathcal{A}$  and  $\mathcal{B}$  are algebraic, it follows that  $A \subset C$  and  $A' \subset C'$ . But then  $f \subset h$  by Lemma 2.1. Thus,  $B(\mathcal{A}, \mathcal{B})$  is algebraic. If  $B(\mathcal{A})$  is algebraic and  $A \subset C \subset B \subset D$  in  $\mathcal{E}(\mathcal{A})$ , then  $\text{Id}_A \sim \text{Id}_B \subset \text{Id}_D$ , hence,  $\text{Id}_A \subset \text{Id}_D$ , whence,  $\text{Id}_{A \cup D} = \text{Id}_{A \cup D}$  belongs to  $B(\mathcal{A})$  – whence,  $A \subset D$ , and it follows that  $\mathcal{A}$  is algebraic.

A test space  $\mathcal{A}$  is **coherent** [3, 4] iff for all events  $A$  and  $B$  of  $\mathcal{A}$ ,  $A \subseteq B^\perp \Rightarrow A \perp B$ .

**2.3 Theorem:** Let  $\mathcal{A}$  and  $\mathcal{B}$  be coherent and  $\kappa$ -uniform. Then  $B(\mathcal{A}, \mathcal{B})$  is also coherent.

Proof: Suppose  $f, g \in \mathcal{E}(B(\mathcal{A}, \mathcal{B}))$  with  $f : A \rightarrow A'$  and  $g : B \rightarrow B'$ . Suppose  $f \subseteq g^\perp$ . Then for every  $x \in A$ ,  $(x, f(x)) \perp (y, g(y))$  for every  $y \in B$ ; hence,  $x \in B^\perp$  and (since  $g$  is surjective),  $f(x) \in B'^\perp$ . Thus,  $A \subseteq B^\perp$  and (since  $f$  is surjective)  $A' \subseteq B'^\perp$ . Since  $\mathcal{A}$  and  $\mathcal{B}$  are coherent,  $A \perp B$  and  $A' \perp B'$ . Thus,  $f \cap g = \emptyset$  and  $f \cup g \in B(\mathcal{E}(\mathcal{A}))$ . Since  $\mathcal{A}$  and  $\mathcal{B}$  are  $n$ -uniform,  $f \perp g$ . Thus,  $B(\mathcal{A}, \mathcal{B})$  is coherent.

A *support* of a test space  $\mathcal{A}$  is a set  $S \subseteq X = \bigcup \mathcal{A}$  such that for all  $E, F \in \mathcal{A}$ ,

$$E \cap S \subseteq F \Rightarrow F \cap S \subseteq E.$$

The usual heuristic is that  $S$  is the set of outcomes that are *possible* in some state of affairs. By way of example, if  $\omega$  is a (probabilistic) state on  $\mathcal{A}$ , then  $S_\omega = \{x \in X \mid \omega(x) > 0\}$  is a support of  $\mathcal{A}$ . Notice that  $X$  is a support, since test spaces are irredundant. It is straight-forward that the union of any collection of supports is a support; hence, the set of all supports of  $\mathcal{A}$  is a complete lattice under set inclusion. More details and motivation will be found in [4].

Let  $\cup B(\mathcal{A}, \mathcal{B}) = Z \subseteq X \times Y$ . Suppose  $S$  and  $T$  are supports of  $\mathcal{A}$ . Then we define

$$S \odot T := [X \times T \cup S \times Y] \cap Z.$$

**2.4 Lemma:** *If  $S$  and  $T$  are supports of  $\mathcal{A}$  and  $\mathcal{B}$ , respectively, then  $S \odot T$  is a support of  $B(\mathcal{A}, \mathcal{B})$ .*

Proof: Suppose  $f : E \rightarrow E'$  and  $g : F \rightarrow F'$  are operations in  $B(\mathcal{A}, \mathcal{B})$ , and that

$$f \cap (S \odot T) = \{(x, f(x)) \mid x \in E \cap S \text{ or } f(x) \in E' \cap T\} \subseteq g.$$

Then  $E \cap S \subseteq F = \text{dom}(g)$  and  $E' \cap T \subseteq F' = \text{ran}(g)$ , whence, as  $S$  and  $T$  are supports,  $E \cap S = F \cap S$  and  $E' \cap T = F' \cap T$ . Moreover,  $f|_{E \cap S} = g|_{F \cap S}$  and  $f^{-1}|_{E' \cap T} = g^{-1}|_{F' \cap T}$ . Hence,  $g \cap (S \odot T) = f \cap (S \odot T)$ . Thus,  $S \odot T$  is a support of  $B(\mathcal{A})$ .

**Remark:** If  $\mu$  is a state on  $\mathcal{A}$ , then  $\mu \circ \pi_1$  is a state on  $B(\mathcal{A}, \mathcal{B})$  (provided that the latter test space exists). Hence, given a state  $\mu$  on  $\mathcal{A}$  and a state  $\nu$  on  $\mathcal{B}$ , we may form a state

$$\mu \odot \nu := \frac{1}{2}(\mu \circ \pi_1 + \nu \circ \pi_2)$$

on  $B(\mathcal{A}, \mathcal{B})$ . It is easily checked that  $S_{\mu \odot \nu} = S_\mu \odot S_\nu$ .

A test space  $\mathcal{A}$  is **regular** iff, for every  $x \in X = \cup \mathcal{A}$ ,  $X \setminus x^\perp$  is a support of  $\mathcal{A}$  [4]. We have:

**2.5 Theorem:** *If  $\mathcal{A}$  and  $\mathcal{B}$  are regular, so is  $B(\mathcal{A}, \mathcal{B})$ .*

Proof: For a typical outcome  $(x, y) \in Z$ , we have

$$Z \setminus (x, y)^\perp = [(X \setminus x^\perp) \times Y \cup X \times (Y \setminus y^\perp)] \cap Z = (X \setminus x^\perp) \odot (X \setminus y^\perp).$$

Since  $\mathcal{A}$  and  $\mathcal{B}$  are regular, this last is a support by Lemma 3. Hence,  $B(\mathcal{A}, \mathcal{B})$  is regular.

Let us adopt the following notation: If  $S$  is a support of a test-space  $\mathcal{A}$  and  $\alpha : E \rightarrow V$  is a  $V$ -valued observable, then we write  $\{\alpha \in A\}$  for the collection of all supports of  $\mathcal{A}$  such that  $\alpha(S \cap E) \subseteq A$ . That is:  $\{\alpha \in A\}$  is the set of all supports making the event  $\alpha^{-1}(A)$  certain to occur if the test  $E$  is made.

**2.6 Lemma:** *Let  $\alpha$  and  $\beta$  be  $V$ -valued questions and  $A \subseteq V$ . Then*

$$\{\alpha \cdot \beta \in A\} = \{\alpha \in A\} \odot \{\beta \in A\}.$$



Proof: Suppose  $\alpha : E \rightarrow V$  and  $\beta : F \rightarrow V$ . Let  $f = \beta^{-1}\alpha = \{(x, y) \in E \times F \mid \alpha(x) = \beta(y)\}$ . Then

$$(S \odot T) \cap f = \{(x, y) \mid \alpha(x) = \beta(y) \ \& \ x \in E \cap S \text{ or } y \in F \cap T\}.$$

Hence,  $S \odot T \cap f \subseteq (\alpha \cdot \beta)^{-1}(A)$  iff  $\alpha(S \cap E) \subseteq A$  and  $\beta(T \cap F) \subseteq A$ .

As a special case of the foregoing, note that  $\alpha \cdot \beta$  is certain to take a value in  $A \subseteq V$  in a state of affairs represented by  $S \odot S$  iff both  $\alpha$  and  $\beta$  are certain to lie in  $A$  in the state of affairs represented by  $S$ .

### 3 The Logic of $B(\mathcal{A}, \mathcal{B})$

If  $\mathcal{A}$  is algebraic, the relation  $\sim$  of perspectivity is an equivalence relation on the set of events of  $\mathcal{A}$ . The set of equivalence classes of events is the **logic** of  $\mathcal{A}$ , here denoted by  $L(\mathcal{A})$ . The equivalence class  $p(A) := \{B \in \mathcal{E}(\mathcal{A}) \mid B \sim A\}$  of an event  $A$  is called the *operational proposition* corresponding to  $A$ . As is well-known,  $L(\mathcal{A})$  can be organized into an orthoalgebra via the partial binary operation  $p(A) \oplus p(B) := p(A \cup B)$ , (well)-defined for pairs of events  $A, B$  with  $A \perp B$ . (For details, see [2] and [3], or [4].)

If  $\mathcal{A}$  and  $\mathcal{B}$  are algebraic, then  $B(\mathcal{A}, \mathcal{B})$  is also algebraic, by Theorem 2.2. In this section, we characterize  $\Pi(B(\mathcal{A}, \mathcal{B}))$  in terms of  $\Pi(\mathcal{A})$  and  $\Pi(\mathcal{B})$  for a large class of algebraic test spaces.

**3.1 Definition:** Events  $A \in \mathcal{E}(\mathcal{A})$  and  $B \in \mathcal{E}(\mathcal{B})$  are **comparable** iff there exists a bijection  $f \in \mathcal{E}(B(\mathcal{A}, \mathcal{B}))$  with  $f : A \rightarrow B$ .

Note that if  $\mathcal{A}$  is  $\kappa$ -uniform, then any two proper events  $A$  and  $B$  of a given cardinality are comparable.

Let  $A$  and  $B$  be comparable events. By Lemma 2.1, the proposition  $p(f)$  corresponding to any (hence, all) bijections  $f : A \rightarrow B$  consists exactly of the union of the sets  $B(C, D)$  of bijections between  $C$  and  $D$  with  $C \sim A$  and  $D \sim B$ . Thus, the proposition  $p(f)$  is completely determined by the pair  $p(A)$  and  $p(B)$ . Let us write  $p(A, B)$  for this proposition.

**3.2 Lemma:** Let  $A, B \in \mathcal{E}(\mathcal{A})$  and  $C, D \in \mathcal{E}(\mathcal{B})$  with  $A$  and  $C$  comparable and  $B$  and  $D$  comparable. If  $p(A, B) \perp p(C, D)$ , then  $A \perp C$ ,  $B \perp D$ ,  $A \cup C$  and  $B \cup D$  are comparable, and  $p(A, B) \oplus p(C, D) = p(A \cup C, B \cup D)$ .

Proof: If  $p(A, B) \perp p(C, D)$  then for every bijection  $f : A \rightarrow B$  and every bijection  $g : C \rightarrow D$ ,  $f \cap g = \emptyset$  and  $f \cup g : A \cup C \rightarrow B \cup D$  belongs to  $\mathcal{E}(B(\mathcal{A}))$  – whence,  $A \perp C$ ,  $B \perp D$ , and  $p(A \cup C, B \cup D) = p(f \cup g) = p(f) \oplus p(g) = p(A, B) \oplus p(C, D)$ .

Note that  $A \perp C, B \perp D$  need not imply that  $p(A, B) \perp p(B, D)$  unless  $\mathcal{A}$  is uniform.

If  $\mathcal{A}$  is uniform, then any two perspective events have the same cardinality. Hence, we may define a map  $\rho : \Pi(\mathcal{A}) \rightarrow \kappa$  (where  $\kappa$  is the cardinality of a test in  $\mathcal{A}$ ) by

$$\rho(p(A)) = \#(A).$$

We call  $\rho(p)$  the **rank** of the proposition  $p \in \Pi(\mathcal{A})$ . Note also that if  $p \perp q$  then  $\rho(p \oplus q) = \rho(p) + \rho(q)$  for all  $p, q \in \Pi(\mathcal{A})$ . The proof of the following is straightforward:

**3.3 Theorem:** *Let  $\mathcal{A}$  and  $\mathcal{B}$  be  $\kappa$ -uniform test spaces with logics  $L$  and  $M$ , respectively. Let*

$$L \times_{\rho} M = \{ (p, q) \in L \times M \mid \rho(p) = \rho(q) \}.$$

*For all  $(p, q), (u, v) \in L \times_{\rho} M$ , write  $(p, q) \perp (u, v)$  iff  $p \perp u$  and  $q \perp v$ , and, if this is the case, set  $(p, q) \oplus (u, v) := (p \oplus q, u \oplus v)$ . Then  $(L \times_{\rho} M, \perp, \oplus)$  is an orthoalgebra, and there is a canonical isomorphism  $L \times_{\rho} M \rightarrow \Pi(B(\mathcal{A}, \mathcal{B}))$  given by  $p(A, B) \mapsto (p(A), p(B))$ .*

**3.4 Definition:** Call an algebraic test space  $\mathcal{A}$  **saturated** iff for every  $A \in \mathcal{E}(\mathcal{A})$  there is some  $x_A \in X = \cup \mathcal{A}$  with  $\{x_A\} \sim A$ .

By way of example, if  $\mathcal{A}$  is any manual, the manual  $\mathcal{A}^{\#}$  of partitions of  $\mathcal{A}$ -operations by  $\mathcal{A}$ -events is saturated, with  $\{\cup A\} \sim A$  for any subset  $A$  of such a partition.

**3.5 Lemma:** *Let  $\mathcal{A}$  and  $\mathcal{B}$  be saturated. Then every bijection between proper events of  $\mathcal{A}$  and  $\mathcal{B}$  can be extended to an element of  $B(\mathcal{A}, \mathcal{B})$ .*

Proof: If  $A$  and  $B$  are proper events of the same cardinality with  $A \subseteq E \in \mathcal{A}$  and  $B \subseteq F \in \mathcal{B}$ , then there exist outcomes  $x$  and  $y$  with  $\{x\} \sim E \setminus A$  and  $\{y\} \sim F \setminus B$ . Since  $\mathcal{A}$  and  $\mathcal{B}$  are algebraic,  $\{x\} \cup A$  and  $\{y\} \cup B$  are tests in  $\mathcal{A}$  and  $\mathcal{B}$ , respectively, to which we may extend any given bijection  $f : A \rightarrow B$  by setting  $f(x) = y$ .

**3.6 Definition:** Let  $L$  and  $M$  be two orthoalgebras. Let

$$L * M := \{ (p, q) \in L \times M \mid p = 0 \Leftrightarrow q = 0 \ \& \ p = 1 \Leftrightarrow q = 1 \}.$$

For  $(p, q)$  and  $(r, s)$  in  $L * M$ , set  $(p, q) \perp (r, s)$  iff  $p \perp r, q \perp s$ , and  $(p \oplus q, r \oplus s) \in L * M$ . If this is the case, define  $(p, q) \oplus (r, s) := (p \oplus r, q \oplus s)$ .

It is easily verified that  $(L * M, \perp, \oplus, (1, 1))$  is an orthoalgebra in which the orthocomplement of an element  $(p, q)$  is given by  $(p, q)' = (p', q')$ .

**3.7 Proposition:** *Let  $\mathcal{A}_1$  and  $\mathcal{A}_2$  be saturated algebraic test spaces. Then*

$$L(B(\mathcal{A}_1, \mathcal{A}_2)) \simeq L(\mathcal{A}_1) * L(\mathcal{A}_2).$$

Proof: Let  $L = L(B(\mathcal{A}_1, \mathcal{A}_2))$  and  $L_i = L(\mathcal{A}_i), i = 1, 2$ . The two coordinate projections  $\pi_i : B(\mathcal{A}_1, \mathcal{A}_2) \rightarrow \mathcal{A}_i$  introduced in the proof of Theorem 1.7 lift to orthoalgebra homomorphisms  $L \rightarrow L_i$ . Since  $\mathcal{A}_1$  and  $\mathcal{A}_2$  are saturated, these are surjections, by Lemma 3.5 above. Hence, we have a natural map  $\phi : L \rightarrow L_1 \times L_2$  given by  $\phi(p) = (\pi_1(p), \pi_2(p))$  for all  $p \in L$ . If  $\phi(p) = (1, q) \in L_1 \times L_2$ , then  $p = p(E, B)$  for some  $E \in \mathcal{A}_1$  and some event  $B \in \mathcal{A}_2$

with  $q = p(B)$ . In order for  $A$  and  $B$  to be comparable, there must exist a bijection  $f \in B(\mathcal{A}_1, \mathcal{A}_2)$  with  $B = f(A)$ . But then  $B \in \mathcal{A}_2$ , whence,  $q = 1$ . Similarly, if  $p = 0$ ,  $q = 0$  in order to preserve comparability. On the other hand, if  $p \in L(\mathcal{A}_1)$ ,  $q \in L(\mathcal{A}_2)$ , and neither  $p$  nor  $q$  is 0 or 1, then, since each manual is saturated, we may choose outcomes  $x \in X_1 = \bigcup \mathcal{A}_1$  and  $y \in X_2 = \bigcup \mathcal{A}_2$  with  $p = p(x)$  and  $q = p(y)$ . Likewise,  $p' = p(x')$  and  $q' = p(y')$  for some outcomes  $x' \in X_1$  and  $y' \in X_2$ . Thus  $\{x, x'\}$  and  $\{y, y'\}$  are two-element tests in  $\mathcal{A}_1$  and  $\mathcal{A}_2$ , respectively, whence,  $f = \{(x, y), (x', y')\}$  belongs to  $B(\mathcal{A}_1, \mathcal{A}_2)$ . Thus,  $\pi(p(f)) = \pi(p(x, y)) = (p, q)$  is defined. The image of  $\phi$  is therefore precisely  $L_1 * L_2$ . It remains to see that  $\phi$  is an faithful (hence, injective) orthoalgebra homomorphism. But this follows from Lemma 3.2.

A **partition of unity** in an orthoalgebra  $L$  is a finite set  $E = \{p_1, \dots, p_n\} \subseteq L \setminus \{0\}$  such that  $p_1 \oplus \dots \oplus p_n = 1$ . The collection  $\mathcal{A}_L$  of all such partitions of unity is easily seen to be a saturated manual, the logic of which is canonically isomorphic to  $L$ .

**3.8 Corollary:** For any orthoalgebras  $L$  and  $M$ ,  $L * M \simeq L(B(\mathcal{A}_L, \mathcal{A}_M))$ .

Call two test spaces  $\mathcal{A}$  and  $\mathcal{B}$  *uniformly compatible* iff every bijection between events of  $\mathcal{A}$  and  $\mathcal{B}$  extends to an element of  $B(\mathcal{A}, \mathcal{B})$ . (By way of example: Any two saturated algebraic test spaces, or any two uniform test spaces). The following generalization of Theorem 3.7 is straightforward. We omit the proof.

**3.9 Proposition:** Let  $\mathcal{A}$  and  $\mathcal{B}$  be uniformly compatible test spaces. There is a canonical embedding of  $L(B(\mathcal{A}, \mathcal{B}))$  into  $L(\mathcal{A}) * L(\mathcal{B})$  given by

$$(p(A, B)) \mapsto (p(A), p(B))$$

for compatible events  $A$  and  $B$ .

## References

- [1] D. Aerts, Foundations of Physics **24** (1994) 1227.
- [2] D. J. Foulis, R. Greechie and G. T. Rüttimann, International Journal of Theoretical Physics **31** (1992) 789-807.
- [3] D. J. Foulis, R. Greechie and G.T. Rüttimann, International Journal of Theoretical Physics **32** (1993) 1675-1689.
- [4] D. J. Foulis, C. Piron and C. H. Randall, Foundations of Physics **13** (1983) 813-842.
- [5] A. M. Gleason, Journal of Mathematics and Mechanics, **6** (1957) 885-893
- [6] S. Gudder, Annales de l'Institut Henri Poincaré, **45** (1986) 311-326.
- [7] G. Kalmbach, **Orthomodular Lattices**, Academic Press, New York, 1983.
- [8] C. Piron, **Foundations of Quantum Physics**, W. A. Benjamin, Reading, 1976.